TWO-PERSON COOPERATIVE GAMES

Part II: The Axiomatic Approach

J. Nash – 1953

The Approach

"Rather than solve the two-person cooperative game by analyzing the bargaining process, one can attack the problem axiomatically by stating general properties that 'any reasonable solution' should possess. By specifying enough such properties one excludes all but one solution."

Notations

- □ Two players: $N = \{1, 2\}$, often denoted by *i* or *j*.
- D is the event "disagreement"
- $\square d = (u_1(D), u_2(D))$
- \square The value to the two players are $\psi_1(\text{B,d})$ and $\psi_2(\text{B,d})$

Axiom I: (weak) Pareto efficiency

 The solution is not weakly dominated by any point in B except itself.

If $(\upsilon_1, \upsilon_2) \in B$ and $\upsilon_1 \ge \psi_1$ and $\upsilon_2 \ge \psi_2$ then $(\upsilon_1, \upsilon_2) = (\psi_1, \psi_2)$

□ The players never agree on a payoff $(\psi_{1,}\psi_{2})$ when there is a preferable payoff <u>for both</u>.

Axiom I: (weak) Pareto efficiency

Roth (1977) "The requirement that a solution be Pareto optimal may be thought of as a requirement of <u>collective rationality</u>, since it states that an outcome may not be chosen if there is another outcome which both player agree is preferable."

Axiom II: Invariance to equivalent utility representations.

- Suppose that the bargaining problem (B',d') is obtained from (B,d) by the transformations *u_i* → *a_iu_i* + *b_i*, for *i* = 1,2 and *a_i* > 0 , which means *d'_i* = *a_id_i* + *b_i B'* = {(*a₁u₁* + *b₁, <i>a₂u₂* + *b₂*): (*u₁, u₂*) ∈ *B*}
 Then ψ_i(B',d')=a_i ψ_i(B,d)+b_i
- Just the numerical values will be changed.

Axiom III: Symmetry

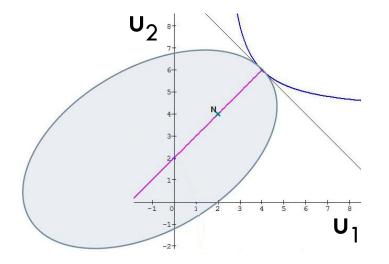
- The solution does not depend on which player is called player one.
- Consider the two bargaining situations (B,d) and (B',d') such that
 - $\square d'_1 = d_2 \text{ and } d_1 = d'_2$
 - **D** $B' = \{(v_1, v_2): (v_2, v_1) \in B\}$

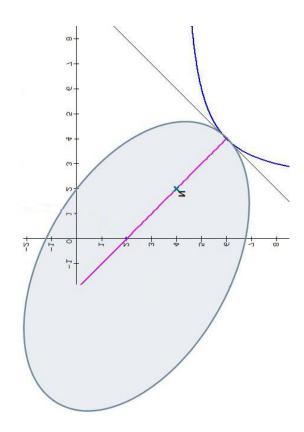
Then the solution point is

 $\psi_1(B',d') = \psi_2(B,d)$ and $\psi_2(B',d') = \psi_1(B,d)$

The difference between the players are the set of strategies and the utility functions.

Axiom III: Symmetry



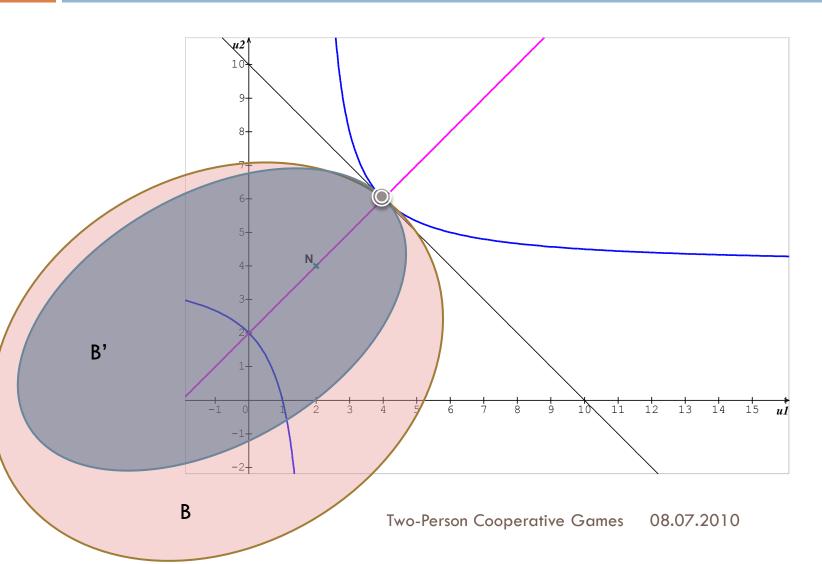


Axiom IV: Independence of irrelevant alternatives

□ If (B,d) and (B',d) are two bargaining situations with B'⊂B and ψ(B,d)∈B'

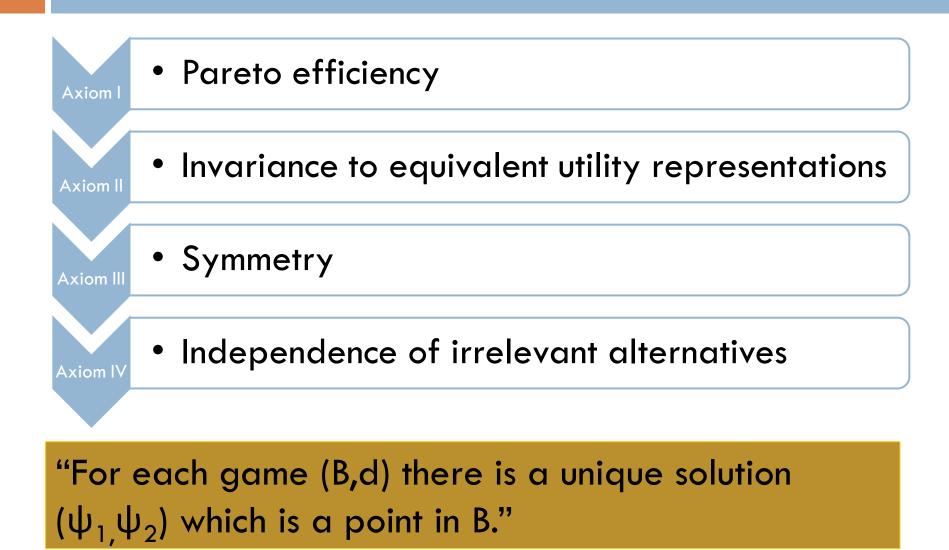
□ Then ψ (B,d)= ψ (B',d)

Axiom IV: Independence of irrelevant alternatives



Discus.

Solution



Solution

The Nash solution is

$$(d_1, d_2) \coloneqq \arg \max_{\substack{(u_1, u_2) \in B \\ u_{iN} = \text{const.} \\ u_i \ge u_{iN}}} (u_1 - u_{1N}) (u_2 - u_{2N})$$

The solution is the utility pair that maximizes the product of the players' utilities (to bargain).

Discussion: Axiom I (Pareto efficiency)

- What if this axiom lacks?
- □ Consider the solution $\psi(B,d)=d$
 - Invariance to equivalent utility representations?
 - Symmetry?
 - Independence of irrelevant alternatives?
- □ This shows how necessary is this axiom

Discussion: Axiom IV (Independence of irrelevant alternatives)

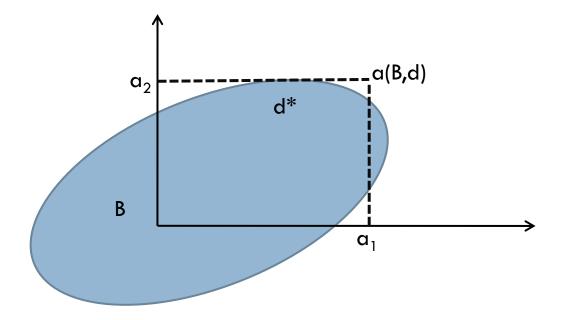
- All points in B (but the solution point ψ) are irrelevant.
- Exclude situations in which the fact that some available agreements influence the outcome.
- "Note that the axiom is satisfied, in particular, by any solution that is defined to be a member of B that maximizes the value of some function." (Rubinstein)

Discussion: Axiom IV (suite)

Nash argues geometrically: "This axiom is equivalent to an axiom of 'localization' of the dependence of the solution point on the shape of the set B. The location of the solution point on the upper-right boundary of B is determined only by the shape of any small segment of the boundary that extends to both sides of it. It does not depend on the rest of the boundary curve."

Idealpunkt a_i

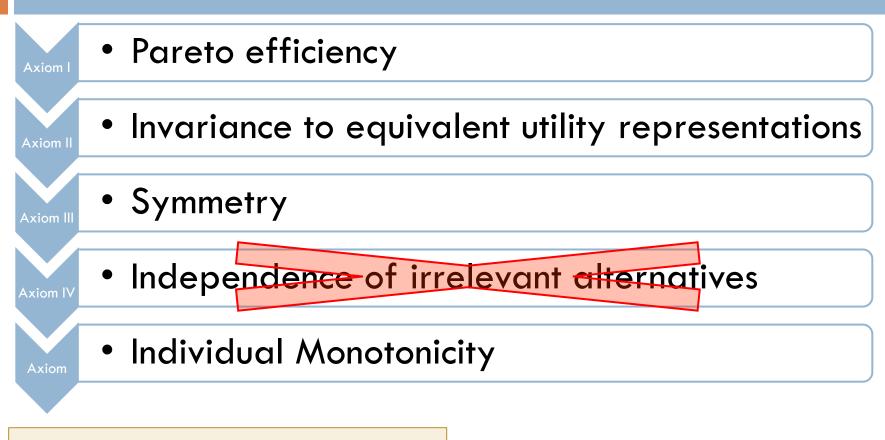
 $a_i(B,d) = \max\{u_i \in R : \text{Es gibt einen Nutzenvektor } u \text{ in } B \text{ mit } u \ge d\}$



An alternative (among others) Axiom: Individual Monotonicity

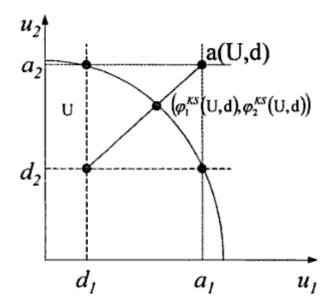
- □ For any (B',d) and (B,d) with B'⊂B and $a_i = a'_i$ we have ψ_i (B',d) ≤ ψ_i (B,d) for i=1,2
- □ For any (B',d) and (B,d) with B'⊂B and $a_i = a'_i$ we have ψ_i (B',d) ≤ ψ_i (B,d) for $i \neq j$

An alternative



Kalai-Smorodinsky's solution

Kalai-Smorodinsky's solution



H. Wiese – Kooperative Spieltheorie, S. 292