

# Applied Cooperative Game Theory

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# Basic example

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- Man owns estate worth 200 euro
- Man dies, leaving debts of 100, 200, 300 euro
- How should the estate be divided?
- First assumptions:
  - The whole estate should be divided
  - Nobody gets more than his claim
  - Everybody gets at least 0.

# Classical Solutions

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- Aristotle proposes  $(100/3, 200/3, 100)$  if the estate is worth 200 euro and  $(50, 100, 150)$  if the estate is worth 300 euro
- Talmud proposes  $(50, 75, 75)$  if the estate is worth 200 euro and  $(50, 100, 150)$  if the estate is worth 300 euro

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## Basic model proposed by O'Neill

Barry O'Neill: "A problem of rights arbitration from the Talmud" 1982  
The set of bankruptcy problems is given by

$$B = \left\{ (N, E, (c_i)_{i \in N}) \mid C := \sum_{i \in N} c_i > E \right\},$$

where

- $N$  is the finite player set,
- $E$  is the worth of the estate,
- $c_i$  is the claim of player  $i$ .

A solution to a bankruptcy problem is a mapping  $S : B \rightarrow \mathbb{R}^N$  such that

$$\begin{aligned} 0 &\leq S_i(N, E, (c_j)_{j \in N}) \leq c_i, i \in N \\ \sum_{j \in N} S_j(N, E, (c_k)_{k \in N}) &= E. \end{aligned}$$

# Axioms #1

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## Definition (Equal Treatment of Equals)

For every bankruptcy problem  $(N, E, (c_k)_{k \in N})$  : If  $c_i = c_j, i, j \in N, i \neq j$ , we then have  $S_i(N, E, (c_k)_{k \in N}) = S_j(N, E, (c_k)_{k \in N})$ .

Similar to Symmetry in cooperative games.

## Definition (Homogeneity)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$  and every  $a \in \mathbb{R}_{++}$  :  $S_i(N, a \cdot E, (a \cdot c_j)_{j \in N}) = a \cdot S_i(N, E, (c_j)_{j \in N}) \forall i \in N$ .

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## Axioms #2

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### Definition (Consistency)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$ : For every  
 $T \subseteq N, S_i(N \setminus T, E - \sum_{j \in S} S_j(N, E, (c_k)_{k \in N}), (c_j)_{j \in N \setminus T}) =$   
 $S_i(N, E, (c_j)_{j \in N}) \forall i \in N \setminus S.$

If a subset of players leave the bargaining process with their rewards, the remaining player's award does not change.

## Axioms #3- Reasonable Bounds

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D.Dominguez, W.Thomson: "A new solution to the problem of adjudicating conflicting claims" 2006

Definition (Reasonable lower bounds on awards)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$  and all  $i \in N$  :  $S_i(N, E, (c_j)_{j \in N}) \geq \frac{1}{n} \min \{c_i, E\}$ .

Definition (Reasonable lower bounds on loses)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$  and all  $i \in N$  :  $c_i - S_i(N, E, (c_j)_{j \in N}) \geq \frac{1}{n} \min \{c_i, C - E\}$ .

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## Axioms #4- Composition

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### Definition (Composition up)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$  and every  $E < E'$  and all  $i \in N$  :

$$S_i(N, E', (c_j)_{j \in N}) = S_i(N, E, (c_j)_{j \in N}) + S_i(N, E' - E, (c_j - S_j(N, E, (c_k)_{k \in N}))_{j \in N})$$

### Definition (Composition down)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$  and every  $E > E'$  and all  $i \in N$  :

$$S_i(N, E', (c_j)_{j \in N}) = S_i(N, E', (S_i(N, E, (c_j)_{j \in N}))) ,$$

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## Solution #1: The Proportional rule

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The easiest way to solve the problem is given by the proportional rule.

Definition (Proportional Rule)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$ :

$$P_i(N, E, (c_j)_{j \in N}) := \frac{E}{\sum_{j \in N} c_j} c_i \forall i \in N.$$

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## Solution #2: The constrained equal award rule

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Definition (constrained equal award rule)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$

$$A_i(N, E, (c_j)_{j \in N}) : = \min \{c_i, \lambda\} \forall i \in N \text{ such that}$$
$$\sum_{j \in N} \min \{c_j, \lambda\} = E.$$

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## Solution #3: The constrained equal lose rule

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Definition (Constrained equal lose rule)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$

$$L_i(N, E, (c_j)_{j \in N}) : = \max \{0, c_i - \lambda\} \forall i \in N \text{ such that}$$

$$\sum_{j \in N} \max \{0, c_j - \lambda\} = E.$$

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## Solution #4: Talmud rule

R.J. Aumann & M.Maschler: Game Theoretic Analysis of a Bankruptcy Problem from the Talmud, 1985.

Definition (Talmud rule)

For every bankruptcy problem  $(N, E, (c_j)_{j \in N})$

$$T_i(N, E, (c_j)_{j \in N}) : = \begin{cases} \min\{c_i/2, \lambda\} & \text{if } E \leq C/2 \\ \max\{c_i/2, c_i - \mu\} & \text{if } E \geq C/2 \end{cases}$$

such that  $\sum_{i \in N} T_i(N, E, (c_j)_{j \in N}) = E.$

# Uniqueness P,L,A

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## Theorem

*The solutions  $P$ ,  $L$ , and  $A$  are the only ones that satisfy*

- *Equal Treatment of Equals*
- *Homogeneity*
- *Consistency*
- *Composition Up*
- *Composition Down*

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# Uniqueness T

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## Theorem

*The solution  $T$  is the only one that satisfies*

- *Consistency*
- *Reasonable Lower Bound On Awards*
- *Reasonable Lower Bound On Loses*

# The nucleolus of a cooperative game #1

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## Definition (excess demand)

Let  $x \in \mathbb{R}^N$  and  $v \in V(N)$ . The excess demand of a coalition  $S$  is given by  $x(S) := v(S) - \sum_{i \in N} x_i$ .

## Definition (order)

Let  $(x(S))_{S \subseteq N}$  be the vector of excess demands. A non-increasing order of the excess demand  $x$  is a function  $\theta : 2^N \rightarrow \mathbb{R}^{2^N}$  such that

$$\theta_i(x) \geq \theta_j(x) \forall i < j.$$

## Definition (lexicographic order)

Let  $x, y \in \mathbb{R}^{2^N}$  and  $\theta : 2^N \rightarrow \mathbb{R}^{2^N}$  be a non-increasing order. We say  $x$  is lexicographically smaller than  $y$  if there is an index  $k \in N$  such that

$$\begin{aligned} \theta_i(x) &= \theta_i(y) \forall i < k \text{ and} \\ \theta_k(x) &< \theta_k(y). \end{aligned}$$

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# The nucleolus of a cooperative game #2

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## Definition

Let  $v \in V(N)$  be a cooperative game. The nucleolus is the efficient payoff vector which lexicographically minimizes the excess demand, i.e.

$$\text{nucleolus}(N, v) = \text{lex} \min_{\substack{x \in \mathbb{R}^N, \\ \sum_{i \in N} x_i = v(N)}} (x(S))_{S \subseteq N}$$

## Corollary

*If the core of the cooperative game  $(N, v)$  is non-empty, the nucleolus of the game  $(N, v)$  lies within the core.*

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# The bankruptcy game

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## Definition

Let  $(N, E, (c_j)_{j \in N})$  be a bankruptcy problem. The associated bankruptcy game is given by

$$v_{E,c}(S) = \max \left\{ 0, E - \sum_{i \in N \setminus S} c_i \right\}.$$

## Theorem (Aumann & Maschler)

*Let  $(N, E, (c_j)_{j \in N})$  be a bankruptcy problem. The nucleolus of the associated bankruptcy game coincides with the Talmud rule of the bankruptcy problem.*

## Definition (Random Arrival Rule)

Let  $(N, E, (c_j)_{j \in N})$  be a bankruptcy problem. The Shapley Value of the associated bankruptcy game is called Random Arrival Rule.

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