

# Advanced Microeconomics

## Static Bayesian games

Harald Wiese

University of Leipzig

## Part E. Bayesian games and mechanism design

1. Static Bayesian games
2. The revelation principle and mechanism design

# Static Bayesian games

1. Conditional probability
2. Introduction and an example
3. Definitions
4. The Cournot model with one-sided cost uncertainty
5. Revisiting mixed-strategy equilibria
6. The first-price auction

# Conditional probability

## Definition

Given:

- ▶ nonempty set  $M$ ,
- ▶ probability distribution  $prob$  on  $M$ ;
- ▶ events  $A$  and  $B$  .

If  $prob(B) \neq 0$ , the conditional probability of  $A$  given  $B$  is defined by

$$prob(A|B) = \frac{prob(A \cap B)}{prob(B)}.$$

If  $prob(A|B) = prob(A)$ ,  $A$  and  $B$  are independent.

# Conditional probability

## Exercises

### Problem

*Throw a dice. What is the conditional probability of 1, 2 or 3 pips (spots) if the number of pips is odd.*

### Problem

*If events  $A$  and  $B$  are independent what about the probability of  $A \cap B$ ?*

# Static Bayesian games

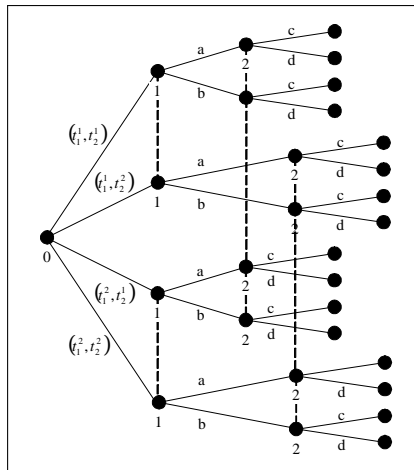
## First-price auction

- ▶ Simultaneous bids where
  - ▶ the highest bidder obtains the object and
  - ▶ pays his bid.
- ▶ Two bidders 1 and 2 who know their own willingness to pay (their “type”).
- ▶ Sequence:
  - ▶ Nature decides the players' types.
  - ▶ Every player learns his own type and can condition his bid on his own type.
- ▶ Static Bayesian games  
= all players act simultaneously after learning their own types.

# Definitions

Static Bayesian game – example for extensive form

- ▶  $N = \{1, 2\}$ ,
- ▶  $T_1 = \{t_1^1, t_1^2\}$  and  $T_2 = \{t_2^1, t_2^2\}$ ,
- ▶  $A_1 = \{a, b\}$  and  $A_2 = \{c, d\}$ .
- ▶ At the initial node, nature chooses a type combination  $t = (t_1, t_2) \in T_1 \times T_2$ .
- ▶  $|T_1|$  information sets for player 1



# Definitions

## Static Bayesian game

### Definition (Static Bayesian game)

A static Bayesian game is a quintuple

$$\Gamma = (N, (A_i)_{i \in N}, (T_i)_{i \in N}, \tau, (u_i)_{i \in N}) = (N, A, T, \tau, u),$$

where

- ▶  $N = \{1, \dots, n\}$  is the player set,
- ▶  $A_i$  is the action set for player  $i \in N$  with Cartesian product  $A = \prod_{i \in N} A_i$  and elements  $a_i$  and  $a$ , respectively,
- ▶  $T = (T_i)_{i \in N}$  is the tuple of type sets  $T_i$  for players  $i \in N$ ,
- ▶  $\tau$  is the probability distribution on  $T$ , and
- ▶  $u_i : A \times T \rightarrow \mathbb{R}$  is player  $i$ 's payoff function (often  $A \times T_i \rightarrow \mathbb{R}$ )



# Definitions

## Beliefs

Ex ante, before the players learn their own types, their beliefs are summarized by  $\tau$ . The (a priori) probability for type  $t_j$  is given by

$$\tau(t_j) := \sum_{t_{-j} \in T_{-j}} \tau(t_{-j}, t_j).$$

## Definition (Belief)

Let  $\Gamma$  be a static Bayesian game with probability distribution  $\tau$  on  $T$ . Player  $i$ 's ex-post (posterior) belief  $\tau_i$  is the probability distribution on  $T_{-i}$  given by the conditional probability

$$\tau_i(t_{-i}) := \tau(t_{-i} | t_j) = \frac{\tau(t_{-i}, t_j)}{\tau(t_j)} = \frac{\tau(t_{-i}, t_j)}{\sum_{t_{-i} \in T_{-i}} \tau(t_{-i}, t_j)}. \quad (1)$$

# Definitions

## Beliefs

### Problem

Two bidders 1 and 2 with  $T_1 = T_2 = \{high, low\}$  (willingness to pay) and

$$\begin{aligned}\tau(high, high) &= \frac{1}{3}, & \tau(high, low) &= \frac{1}{3}, \\ \tau(low, high) &= \frac{1}{9}, & \tau(low, low) &= \frac{2}{9}.\end{aligned}$$

### Find

- ▶  $\tau(t_2 = high)$  (ex ante) and
- ▶  $\tau_1(t_2 = high)$  if player 1 has learned that his own willingness to pay is high (ex post).

# Definitions

## Actions, strategies, and equilibria

### Definition

Let  $\Gamma$  be a static Bayesian game.

- ▶ A strategy for player  $i \in N$  is a function  $s_i : T_i \rightarrow A_i$ . We sometimes write  $s(t)$  instead of  $(s_1(t_1), \dots, s_n(t_n)) \in A$ .
- ▶ A strategy combination  $s^* = (s_1^*, s_2^*, \dots, s_n^*)$  is a Bayesian equilibrium (ex post) if

$$s_i^*(t_i) \in \arg \max_{a_i \in A_i} \sum_{t_{-i} \in T_{-i}} \tau_i(t_{-i}) u_i(a_i, s_{-i}^*(t_{-i}), t_i, t_{-i})$$

holds for all  $i \in N$  and all  $t_i \in T_i$  obeying  $\tau(t_i) > 0$ .

The qualification  $\tau(t_i) > 0$  is necessary because  $\tau_i(t_{-i})$  is ill-defined otherwise.  $\tau(t_i) = 0$  implies “anything goes”.

# The Cournot model with one-sided cost uncertainty

## The model

Static Bayesian game  $\Gamma = (N, (A_1, A_2), (T_1, T_2), \tau, (u_1, u_2))$  with

- ▶ the set of two firms  $N = \{1, 2\}$ ,
- ▶ the action sets  $A_1 = A_2 = [0, \infty)$ ,
- ▶ the type sets  $T_1 = \{c_1^l, c_1^h\} = \{15, 25\}$ ,  $T_2 = \{20\}$ ,  
( $c_2 = 20$  known to both,  $c_1$  known to 1, only)
- ▶ probability distribution  $\tau$  on  $T$  given by  
 $\tau(15, 20) = \tau(25, 20) = \frac{1}{2}$
- ▶ inverse demand function given by  $p(X) = 80 - X$  and hence
- ▶ the payoff functions  
 $u_2(x_1, x_2, t_2) = (p(X) - c_2)x_2 = (80 - (x_1 + x_2) - 20)x_2$   
and

$$u_1(x_1, x_2, t_1) = \begin{cases} (80 - (x_1 + x_2) - 15)x_1, & t_1 = c_1^l \\ (80 - (x_1 + x_2) - 25)x_1, & t_1 = c_1^h \end{cases}$$

# The Cournot model with one-sided cost uncertainty

## The model

The types are independent (trivial).

- ▶ Ex-ante probabilities:

$$\tau(20) : = \tau(15, 20) + \tau(25, 20) = \frac{1}{2} + \frac{1}{2} = 1,$$

$$\tau(15) : = \tau(15, 20) = \frac{1}{2}, \text{ and}$$

$$\tau(25) : = \tau(25, 20) = \frac{1}{2}.$$

- ▶ Ex-post probability for  $c_2 = 20$  (player 1's belief):

$$\tau_1(20) = \frac{\tau(t_1, 20)}{\tau(t_1)} = \frac{\frac{1}{2}}{\frac{1}{2}} = 1 = \tau(20).$$

- ▶ Ex-post probability for  $c_1 = 15$  (player 2's belief):

$$\tau_2(15) = \frac{\tau(15, 20)}{\tau(20)} = \frac{\frac{1}{2}}{1} = \frac{1}{2} = \tau(15).$$

# The Cournot model with one-sided cost uncertainty

## The static Bayesian equilibrium

- ▶ Strategy sets:
  - ▶ player 2:  $S_2 = \{s_2 : \{c_2\} \rightarrow [0, \infty)\}$  to be identified with  $A_2$
  - ▶ player 1:  $S_1 = \left\{s_1 : \left\{c_1^l, c_1^h\right\} \rightarrow [0, \infty)\right\}$ .
- ▶ Firm 1's choice depends on its type:

$$\begin{aligned} s_1^R(t_1) &= \begin{cases} \arg \max_{x_1 \in [0, \infty)} (80 - (x_1 + x_2) - c_1^l) x_1, & t_1 = c_1^l \\ \arg \max_{x_1 \in [0, \infty)} (80 - (x_1 + x_2) - c_1^h) x_1, & t_1 = c_1^h \end{cases} \\ &= \begin{cases} \frac{65}{2} - \frac{1}{2}x_2, & t_1 = c_1^l \\ \frac{55}{2} - \frac{1}{2}x_2, & t_1 = c_1^h \end{cases} \end{aligned}$$

# The Cournot model with one-sided cost uncertainty

## The static Bayesian equilibrium

- ▶ Firm 2's profit is the expected value

$$\begin{aligned} & \tau(15) \left( 80 - [x_1^l + x_2] - 20 \right) x_2 \\ & + \tau(25) \left( 80 - [x_1^h + x_2] - 20 \right) x_2 \\ = & \left( 60 - \frac{1}{2} [x_1^l + x_1^h] \right) x_2 - x_2^2 \end{aligned}$$

which leads to the reaction function

$$x_2^R = s_2^R(20) = 30 - \frac{1}{4} [x_1^l + x_1^h].$$

- ▶ Three equations in the three unknowns  $x_2$ ,  $x_1^l$ , and  $x_1^h$ .
- ▶ They lead to the Nash equilibrium

$$x_2^* = 20, s_1^* \left( c_1^l \right) = \frac{45}{2}, \text{ and } s_1^* \left( c_1^h \right) = \frac{35}{2}.$$

# Revisiting mixed-strategy equilibria

## Continuous types

- ▶ We assume two players  $i = 1, 2$  with types  $t_i$  from  $T_i = [0, x]$ ,  $x > 0$ .
- ▶ Types are independent.

- ▶ Density function:

$$\tau^x(a) = \begin{cases} \frac{1}{x}, & a \in [0, x] \\ 0, & a \notin [0, x] \end{cases}$$

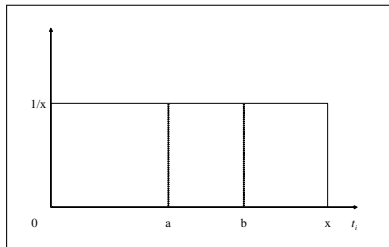
- ▶ Hence, for  $0 \leq a \leq b \leq x$ ,

$$\tau^x([a, b]) = \int_a^b \tau^x(t) dt = \frac{b-a}{x}.$$

- ▶ The probability for a specific type  $a$  is zero:

$$\tau^x([a, a]) = \frac{a-a}{x} = 0.$$

- ▶ Careful: Distinguish  $\tau^x([a, a])$  from  $\tau^x(a)$ .





# Revisiting mixed-strategy equilibria

## Introducing uncertainty

- ▶ Static Bayesian games allow a fresh look on mixed equilibria.
- ▶ For a given matrix game with a mixed-strategy equilibrium, we construct a sequence of static Bayesian games that converges towards that game.
- ▶ In every Bayesian game, no player  $i \in N$  randomizes. However, from the point of view of the other players from  $N \setminus \{i\}$  who do not know the type  $t_i$ , it may well seem as if player  $i$  is a randomizer.

# Revisiting mixed-strategy equilibria

Introducing uncertainty

		Peter	
		theatre	football
Cathy	theatre	$2 + t_C, 1$	$0, 0$
	football	$0, 0$	$1, 2 + t_P$

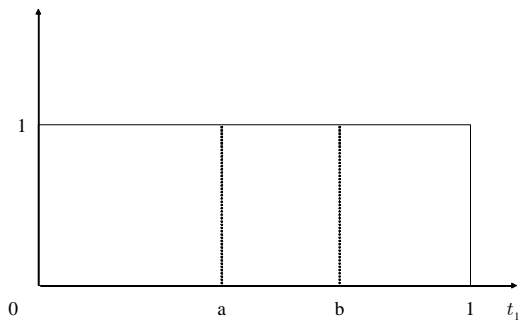
Three equilibria.  $((\frac{2}{3}, \frac{1}{3}), (\frac{1}{3}, \frac{2}{3}))$  is the properly mixed one for  $t_C = t_P = 0$ .

Result (as shown in the manuscript): Cathy's probabilities in the static Bayesian games (which depend on  $x$ ) converge towards her equilibrium mixed strategy as  $x$  goes towards 0.

# The first-price auction

## The model

- ▶ Two bidders 1 and 2 with types  $t_1, t_2 \in [0, 1]$ .
- ▶ Player 1's distribution is given by



# The first-price auction

## The model

Formally, the first-price auction is the static Bayesian game

$$\Gamma = (N, (A_1, A_2), (T_1, T_2), \tau, (u_1, u_2))$$

where

- ▶  $N = \{1, 2\}$  is the set of the two bidders,
- ▶  $A_1 = A_2 = [0, \infty)$  are the sets of bids chosen by the bidders,
- ▶  $T_1 = T_2 = [0, 1]$  are the type sets,
- ▶  $\tau$  is the a probability distribution on  $T$  given by  
 $\tau([a, b], [c, d]) = (b - a)(d - c)$  where  $0 \leq a \leq b \leq 1$  and  $0 \leq c \leq d \leq 1$  hold, and

# The first-price auction

## The model

- ▶ the payoff functions are defined by

$$u_1(a, t_1) = \begin{cases} t_1 - a_1, & a_1 > a_2, \\ \frac{t_1 - a_1}{2}, & a_1 = a_2, \\ 0, & a_1 < a_2, \end{cases} \quad \text{and}$$
$$u_2(a, t_2) = \begin{cases} 0 & a_1 > a_2, \\ \frac{t_2 - a_2}{2} & a_1 = a_2, \\ t_2 - a_2 & a_1 < a_2. \end{cases}$$

# The first-price auction

## Solution

In order to solve the first-price auction, we use the ex-post equilibrium definition. For example, if player 1 is of type  $t_1 \in [0, 1]$ , his condition for the equilibrium strategy combination  $(s_1^*, s_2^*)$  is

$$s_1^*(t_1) \in \arg \max_{a_1 \in A_1} \left( (t_1 - a_1) \underbrace{\tau(\{t_2 \in [0, 1] : a_1 > s_2^*(t_2)\})}_{\substack{\text{probability that player 1's bid} \\ \text{is higher than player 2's bid}}} \right. \\ \left. + \frac{1}{2} (t_1 - a_1) \underbrace{\tau(\{t_2 \in [0, 1] : a_1 = s_2^*(t_2)\})}_{\text{probability for equal bids = 0}} \right).$$

# The first-price auction

## Solution

- ▶ Following Gibbons (1992, pp. 152), we restrict our search for equilibrium strategies to linear strategies of the forms

$$s_1^*(t_1) = c_1 + d_1 t_1 \quad (d_1 > 0),$$

$$s_2^*(t_2) = c_2 + d_2 t_2 \quad (d_2 > 0).$$

- ▶ By

$$\begin{aligned} & \tau(\{t_2 \in [0, 1] : a_1 > c_2 + d_2 t_2\}) \\ &= \tau\left(\left\{t_2 \in [0, 1] : t_2 < \frac{a_1 - c_2}{d_2}\right\}\right) \\ &= \tau\left(\left[0, \frac{a_1 - c_2}{d_2}\right]\right) = \frac{a_1 - c_2}{d_2}, \end{aligned}$$

... (next slide)

# The first-price auction

## Solution

- ▶ player 1's maximization problem is solved by

$$s_1^R(t_1) = \arg \max_{a_1 \in A_1} (t_1 - a_1) \frac{a_1 - c_2}{d_2} = \frac{c_2 + t_1}{2}.$$

- ▶ Player 1's best response to  $c_2$  (and hence to player 2's strategy) is a linear strategy with  $c_1 = \frac{c_2}{2}$  and  $d_1 = \frac{1}{2}$ .
- ▶ Analogously, bidder 2's best response is

$$s_2^R(t_2) = \frac{c_1 + t_2}{2}$$

with  $c_2 = \frac{c_1}{2}$  and  $d_2 = \frac{1}{2}$ .

- ▶ Now,  $c_1 = \frac{c_2}{2} = \frac{\frac{c_1}{2}}{2} = \frac{c_1}{4}$  implies  $c_1 = 0$ . Equilibrium candidate:

$$s_1^* : [0, 1] \rightarrow \mathbb{R}_+, \quad t_1 \mapsto s_1^*(t_1) = \frac{t_1}{2} \text{ and}$$

$$s_2^* : [0, 1] \rightarrow \mathbb{R}_+, \quad t_2 \mapsto s_2^*(t_2) = \frac{t_2}{2}$$



# The first-price auction

## Solution

- ▶ These strategies form an equilibrium because the strategies are best responses to each other. If player 2 uses the (linear) strategy  $s_2^*$ ,  $s_1^*(t_1) = \frac{t_1}{2}$  is a best response as shown above. Thus,  $s_1^*$  turns out to be a linear strategy.
- ▶ Therefore, we have found an equilibrium in linear strategies but cannot exclude the possibility of an equilibrium in non-linear strategies.

# The first-price auction

## First-price or second-price auction?

- ▶ We now take the auctioneer's perspective and ask the question whether the first-price auction is preferable to the second-price auction. The auctioneer compares the prices
  - ▶  $\min(t_1, t_2)$  for the second-price auction and
  - ▶  $\max\left(\frac{1}{2}t_1, \frac{1}{2}t_2\right)$  for the first-price auction.
- ▶ We assume a risk-neutral auctioneer who maximizes the expected payoff.
- ▶ It can be shown (see manuscript) that the auctioneer is indifferent between the first-price and the second-price auction!

## Further exercise

Game  $G_A$  with probability  $p > \frac{1}{2}$ , game  $G_B$  with probability  $1 - p$ .  
Assume  $L > M > 1$ .

$G_A$	left	right	$G_B$	left	right
up	$M, M$	$1, -L$	up	$0, 0$	$1, -L$
down	$-L, 1$	$0, 0$	down	$-L, 1$	$M, M$

- (a) Assume that both players are informed which game they play before they choose their actions. Formulate this game as a static Bayesian game!
- (b\*) Assume that player 1 learns whether they play  $G_A$  or  $G_B$  while player 2 does not. Formulate this game as a static Bayesian game and determine all of its equilibria!